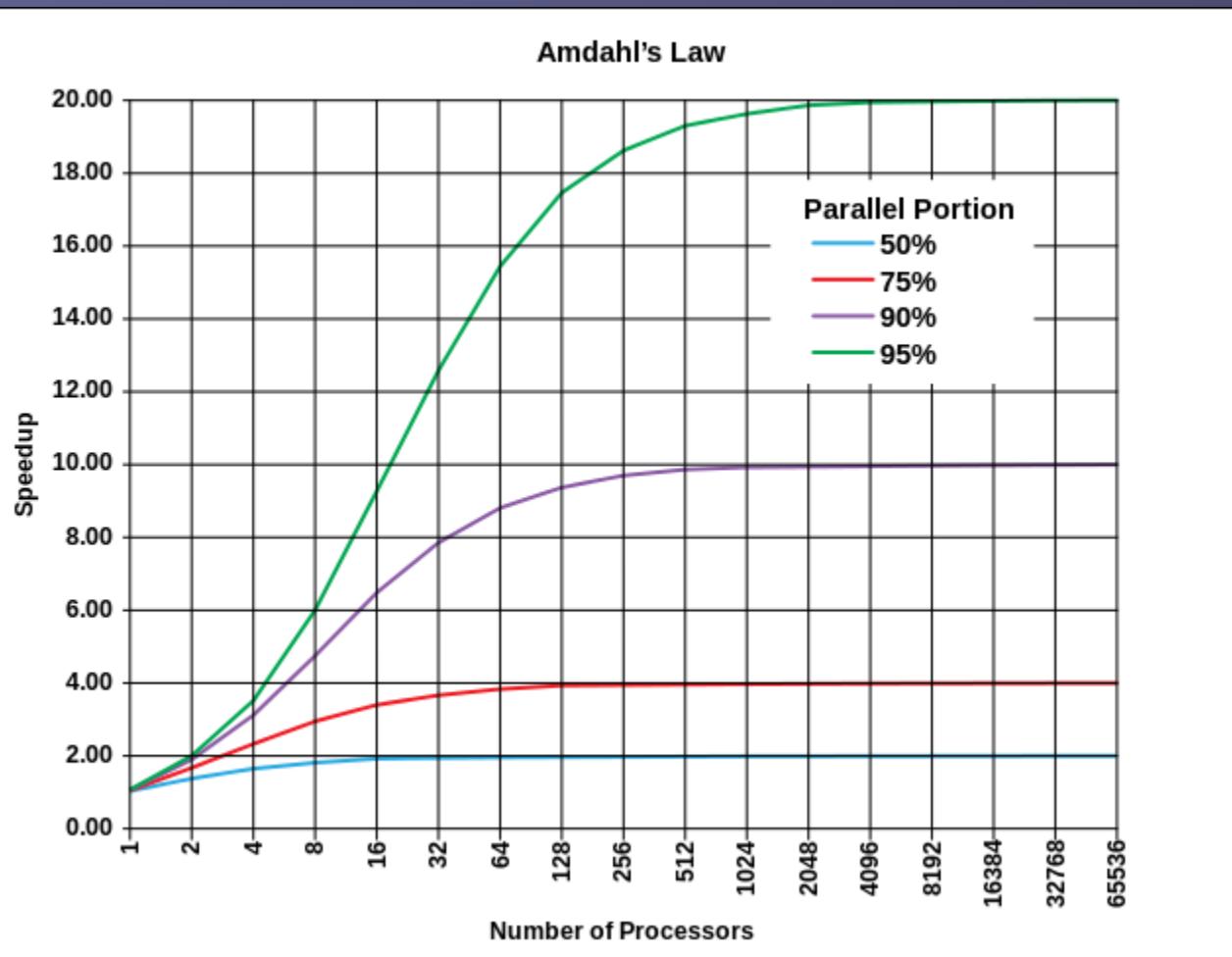


Parallel Computing

Some basic ideas

Amdahl's Law (Gene Amdahl 1967)



Evolution according to Amdahl's law of the theoretical speedup of the execution of a program in function of the number of processors executing it, for different values of p. The speedup is limited by the serial part of the program. For example, if 95% of the program can be parallelized, the theoretical maximum speedup using parallel computing would be 20 times.

Calculate Amdahl's Law:

Let X be the part of my program (in terms of computing time) which can be parallelised. The sequential computing time T_{seq} is normalized to unity (1), and can be expressed as:

$$T_{\text{seq}} = 1 = X + (1-X)$$

The parallel computing time T_{par} under ideal conditions (ideal load balancing, ultrafast communication):

$$T_{\text{par}} = X/p + (1-X)$$

with processor number (core number) p ;

Then the speed-up of the program $S = T_{\text{seq}} / T_{\text{par}}$:

$$S = 1 / (1-X+X/p) \quad ;$$

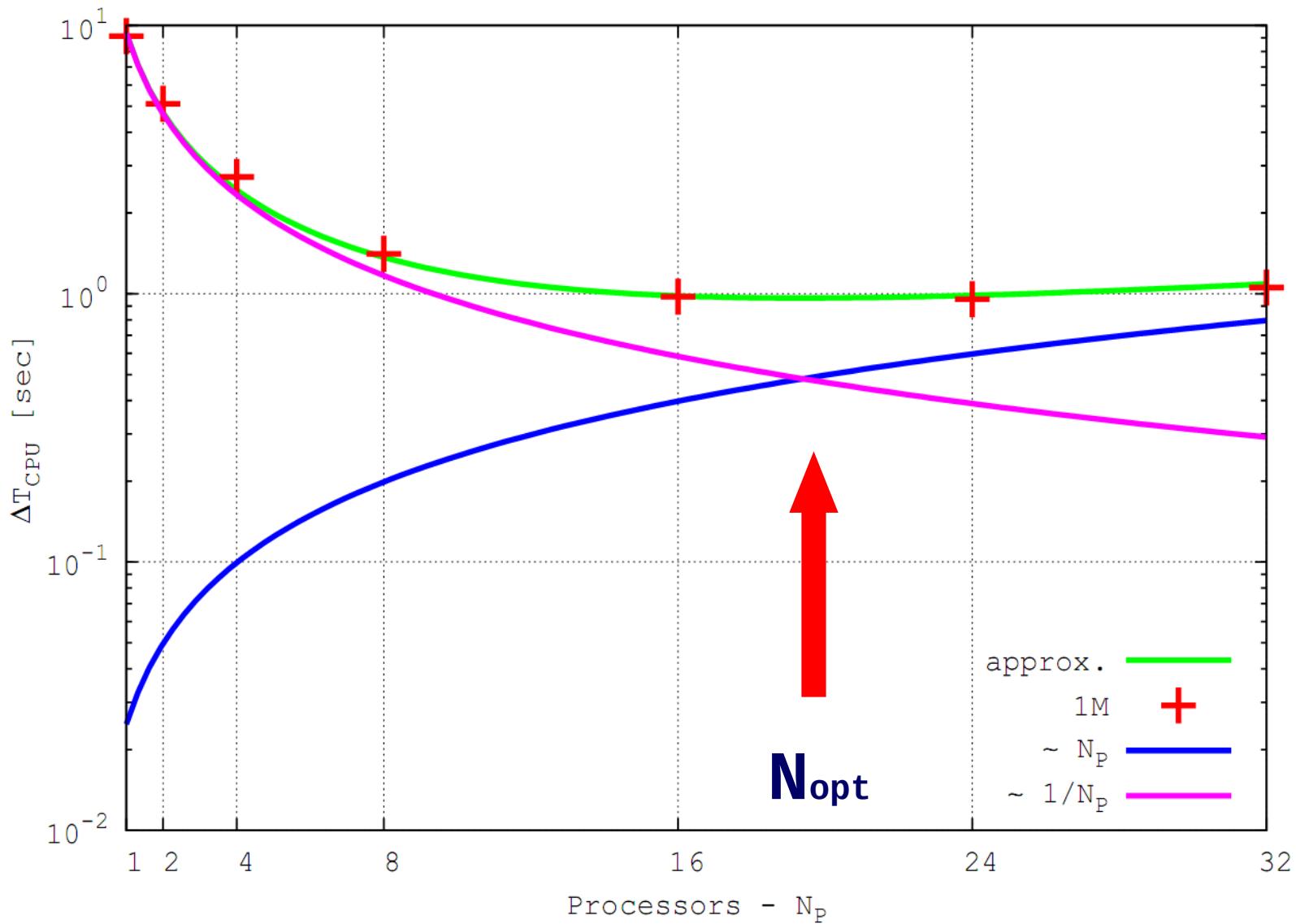
Note: $T_{\text{par}}/T_{\text{seq}} = 1/S$ (sometimes also plotted)

Note the limit of S for large p is very large: $S = 1/(1-X)$. And if $X \sim 1$: $S \sim p$
With communication overhead:

$$T_{\text{par}} = X/p + (1-X) + T_{\text{comm}} \quad \rightarrow \quad S = 1 / (1-X+X/p+T_{\text{comm}})$$

If T_{comm} independent of p we have for large p : $S = 1 / (1-X + T_{\text{comm}}) = \text{const.}$

Parallel code on cluster



Strong and Soft Scaling

- Strong Scaling: Fixed Problem size, increase p
- Soft Scaling: Increase Problem size, increase p
(constant amount of work per processing element)

Ansatz for Soft Scaling:

$$\rightarrow T_{\text{seq}} = p (X + (1-X))$$

$$\rightarrow T_{\text{par}} = X + p (1-X)$$

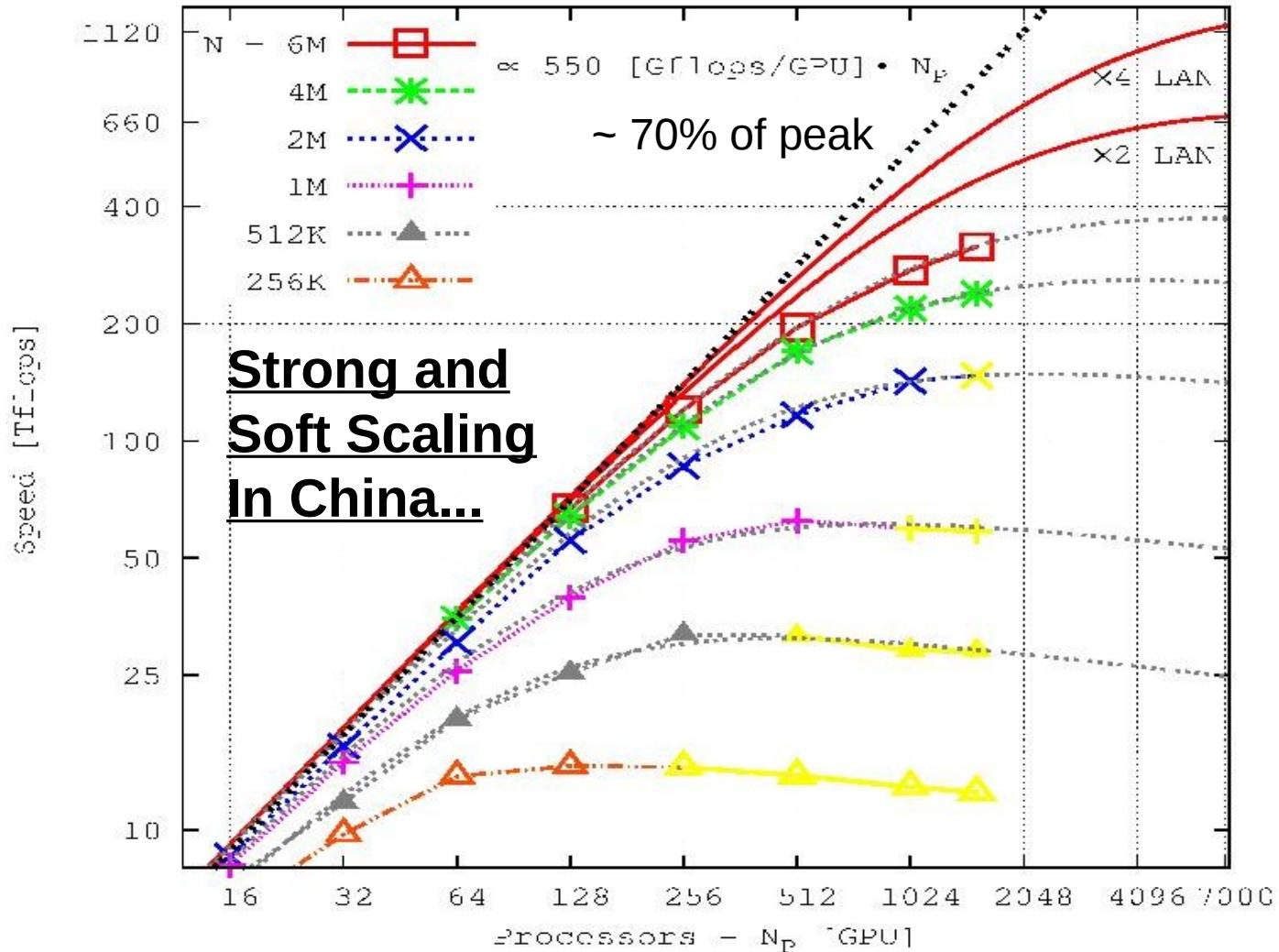
$$\rightarrow S = T_{\text{seq}} / T_{\text{par}} = p / (X + p (1-X))$$

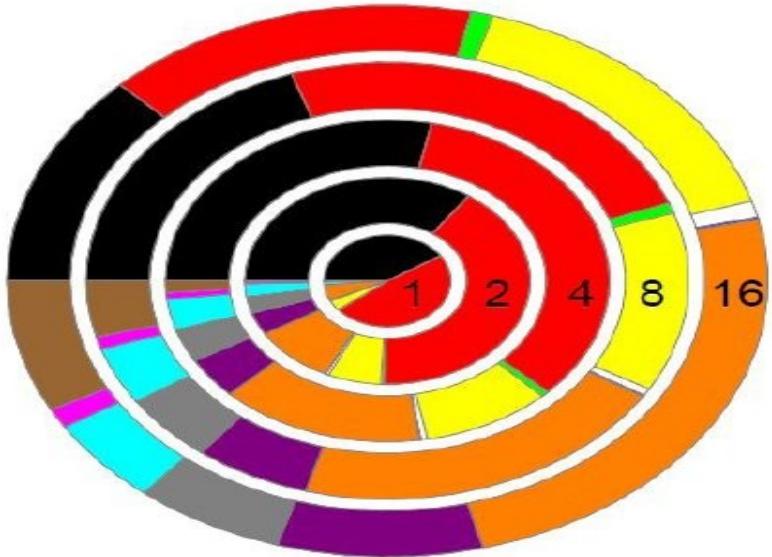
If $X \sim 1$: $S = p$; $T_{\text{par}} = X = \text{const.}$

350 Teraflop/s
1600 GPUs.
440 cores
 $= 704.000$
GPU-Cores

Using
Mole-8.5
of
IPE/CAS
Beijing

Berczik et al.
2013

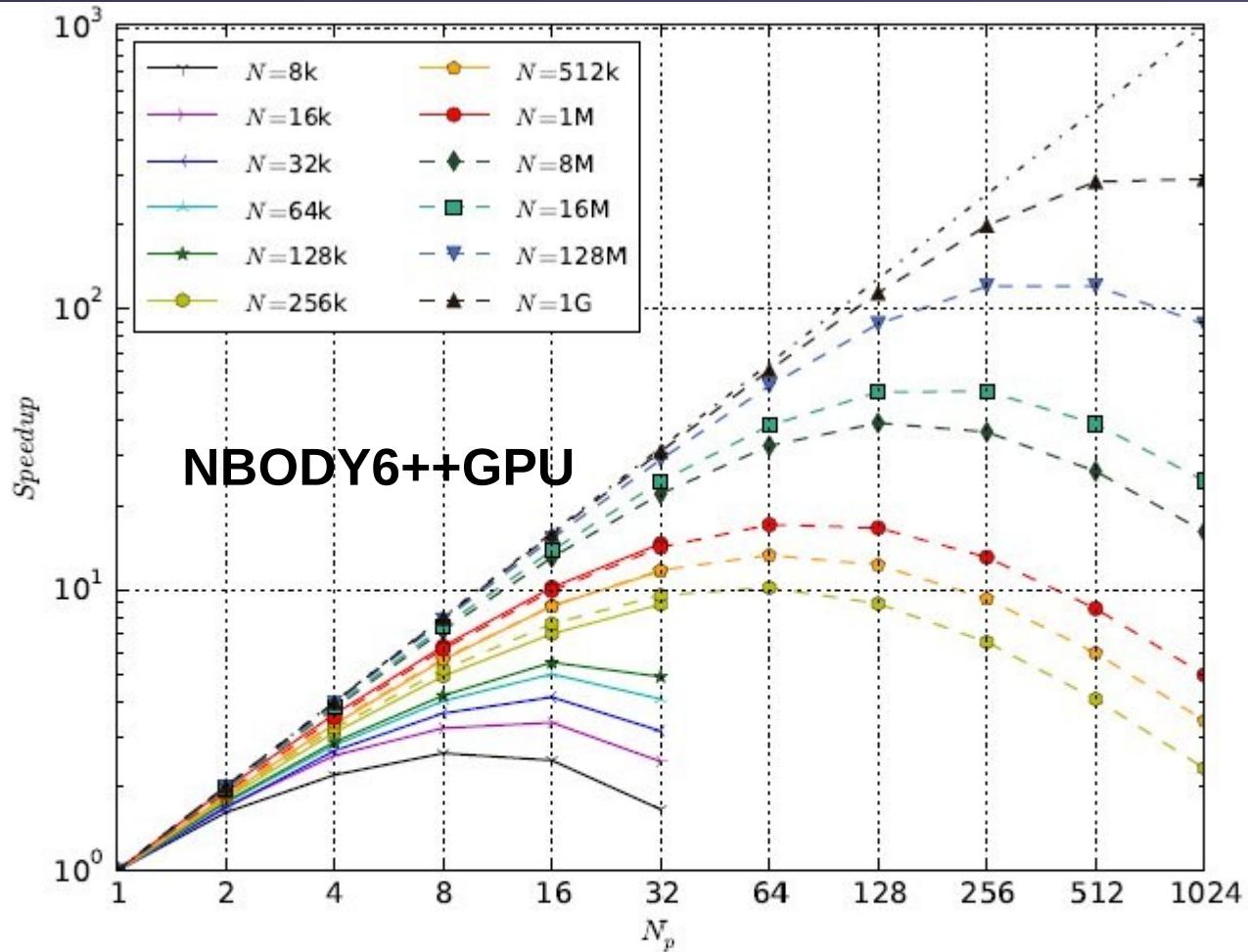




- | | |
|-----------|-----------|
| ■ Reg. | ■ Comm.R. |
| ■ Irr. | ■ Send.I. |
| ■ Pred. | ■ Send.R. |
| ■ Init.B. | ■ Barr. |
| □ Adjust | |
| ■ KS | |
| ■ Move | |
| ■ Comm.I. | |

Table 1 Main components of NBODY6++

Description	Timing variable	Expected scaling		Fitting value [sec]
		N	N_p	
Regular force computation	T_{reg}	$\mathcal{O}(N_{\text{reg}} \cdot N)$	$\mathcal{O}(N_p^{-1})$	$(2.2 \cdot 10^{-9} \cdot N^{2.11} + 10.43) \cdot N_p^{-1}$
Irregular force computation	T_{irr}	$\mathcal{O}(N_{\text{irr}} \cdot \langle N_{nb} \rangle)$	$\mathcal{O}(N_p^{-1})$	$(3.9 \cdot 10^{-7} \cdot N^{1.76} - 16.47) \cdot N_p^{-1}$
Prediction	T_{pre}	$\mathcal{O}(N^{kn_p})$	$\mathcal{O}(N_p^{-kp_p})$	$(1.2 \cdot 10^{-6} \cdot N^{1.51} - 3.58) \cdot N_p^{-0.5}$
Data moving	T_{mov}	$\mathcal{O}(N^{kn_{m1}})$	$\mathcal{O}(1)$	$2.5 \cdot 10^{-6} \cdot N^{1.29} - 0.28$
MPI communication (regular)	T_{mcr}	$\mathcal{O}(N^{kn_{cr}})$	$\mathcal{O}(kp_{cr} \cdot \frac{N_p - 1}{N_p})$	$(3.3 \cdot 10^{-6} \cdot N^{1.18} + 0.12)(1.5 \cdot \frac{N_p - 1}{N_p})$
MPI communication (irregular)	T_{mci}	$\mathcal{O}(N^{kn_{ci}})$	$\mathcal{O}(kp_{ci} \cdot \frac{N_p - 1}{N_p})$	$(3.6 \cdot 10^{-7} \cdot N^{1.40} + 0.56)(1.5 \cdot \frac{N_p - 1}{N_p})$
Synchronization	T_{syn}	$\mathcal{O}(N^{kn_s})$	$\mathcal{O}(N_p^{kp_s})$	$(4.1 \cdot 10^{-8} \cdot N^{1.34} + 0.07) \cdot N_p$
Sequential parts on host	T_{host}	$\mathcal{O}(N^{kn_h})$	$\mathcal{O}(1)$	$4.4 \cdot 10^{-7} \cdot N^{1.49} + 1.23$



Huang, Berczik, Spurzem, Res. Astron. Astroph. 2016, 16, 11.

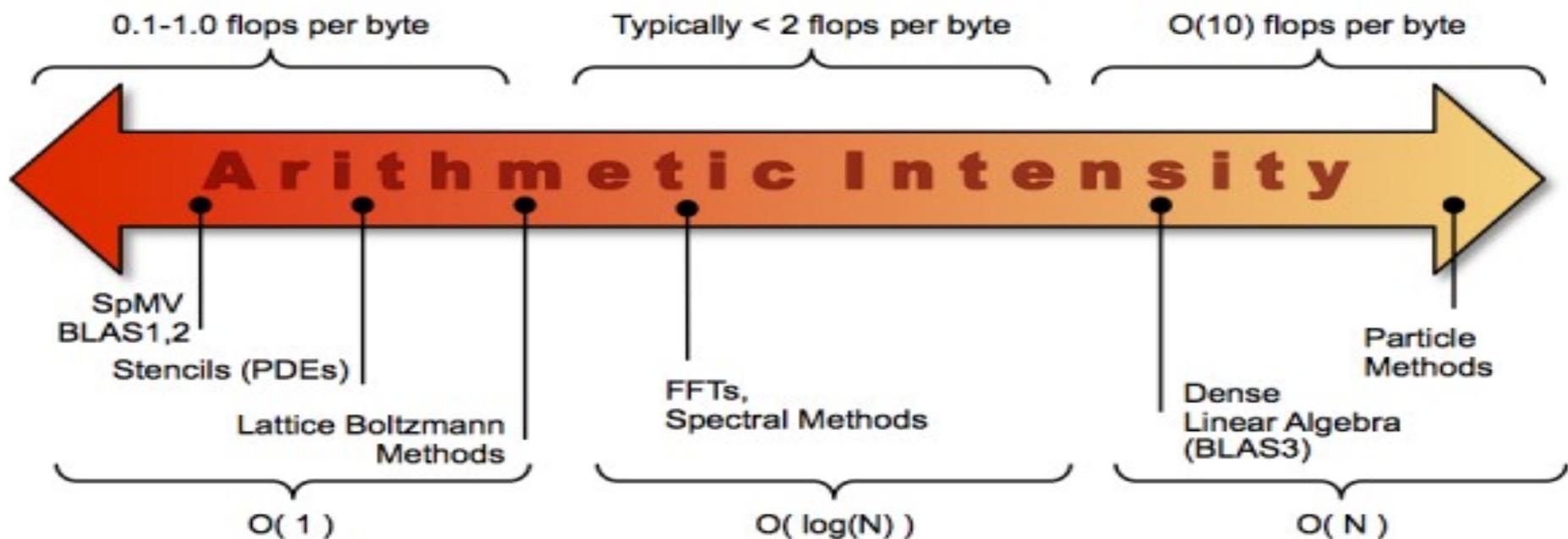
Fig. 2 The speed-up (S) of NBODY6++ as a function of particle number (N) and processor number (N_p). Solid points are the measured speed-up ratio between sequential and parallel wall-clock time, dash lines predict the performance of larger scale simulations further. The symbols used in figure have the magnitudes: $1k = 1,024$, $1M = 1k^2$ and $1G = 1k^3$.

Roofline Performance Model (LBL)

<http://crd.lbl.gov/departments/computer-science/PAR/research/roofline>

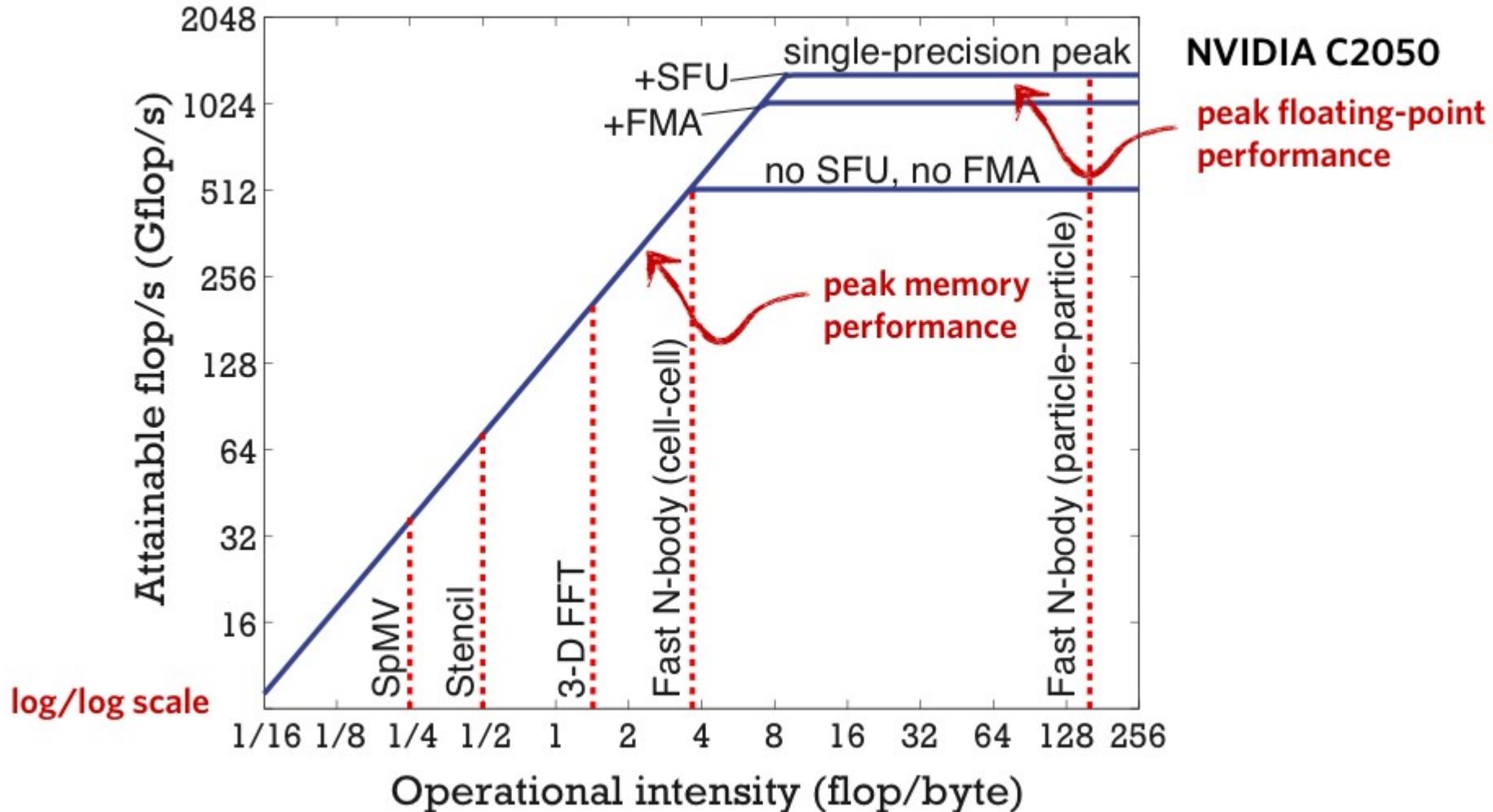
Arithmetic Intensity

The core parameter behind the Roofline model is Arithmetic Intensity. Arithmetic Intensity is the ratio of total floating-point operations to total data movement (bytes).



Roofline Performance Model (LBL)

http://lorenabarba.com/wp-content/uploads/2012/01/roofline_slide.png



Parallel Computing

Timing and Debugging

Wrap-Up of CUDA

Histogram

Matrix Multiplication (expect Friday)

Before we start...

Some nice ideas:

/home/Tit4/lecture60/gpu-course/00_error/

/home/Tit4/lecture60/gpu-course/4_dot/dot_special.cu

Recap of 6: dot_perfect.cu :

Fat Threads! New variable gridDim.x !

Block Reduction on Host instead of AtomicAdd!

Also used for histogram later.

Timing with CUDA Event API

```
int main ()
{
    cudaEvent_t start, stop;
    float time;

    cudaEventCreate (&start);
    cudaEventCreate (&stop);

    cudaEventRecord (start, 0);

    someKernel <<<grids, blocks, 0, 0>>> (...);

    cudaEventRecord (stop, 0);
    cudaEventSynchronize (stop); ← Ensures kernel execution has completed

    cudaEventElapsedTime (&time, start, stop);

    cudaEventDestroy (start);
    cudaEventDestroy (stop);

    printf ("Elapsed time %f sec\n", time*.001);

    return 1;
}
```

CUDA Event API Timer are,

- OS independent
- High resolution
- Useful for timing asynchronous calls

← Ensures kernel execution has completed

Standard CPU timers will not measure the timing information of the device.

CUDA – GNU Debugger – CUDA-gdb

<http://docs.nvidia.com/cuda/cuda-gdb/index.html>

 **DEVELOPER ZONE** **CUDA TOOLKIT DOCUMENTATION**  **Search**

CUDA Toolkit v7.5 CUDA-GDB CUDA-GDB (PDF) - v7.5 (older) - Last updated September 1, 2015 - [Send Feedback](#) -     

CUDA-GDB

1. Introduction

This document introduces CUDA-GDB, the NVIDIA® CUDA® debugger for Linux and Mac OS.

1.1. What is CUDA-GDB?

CUDA-GDB is the NVIDIA tool for debugging CUDA applications running on Linux and Mac. CUDA-GDB is an extension to the x86-64 port of GDB, the GNU Project debugger. The tool provides developers with a mechanism for debugging CUDA applications running on actual hardware. This enables developers to debug applications without the potential variations introduced by simulation and emulation environments.

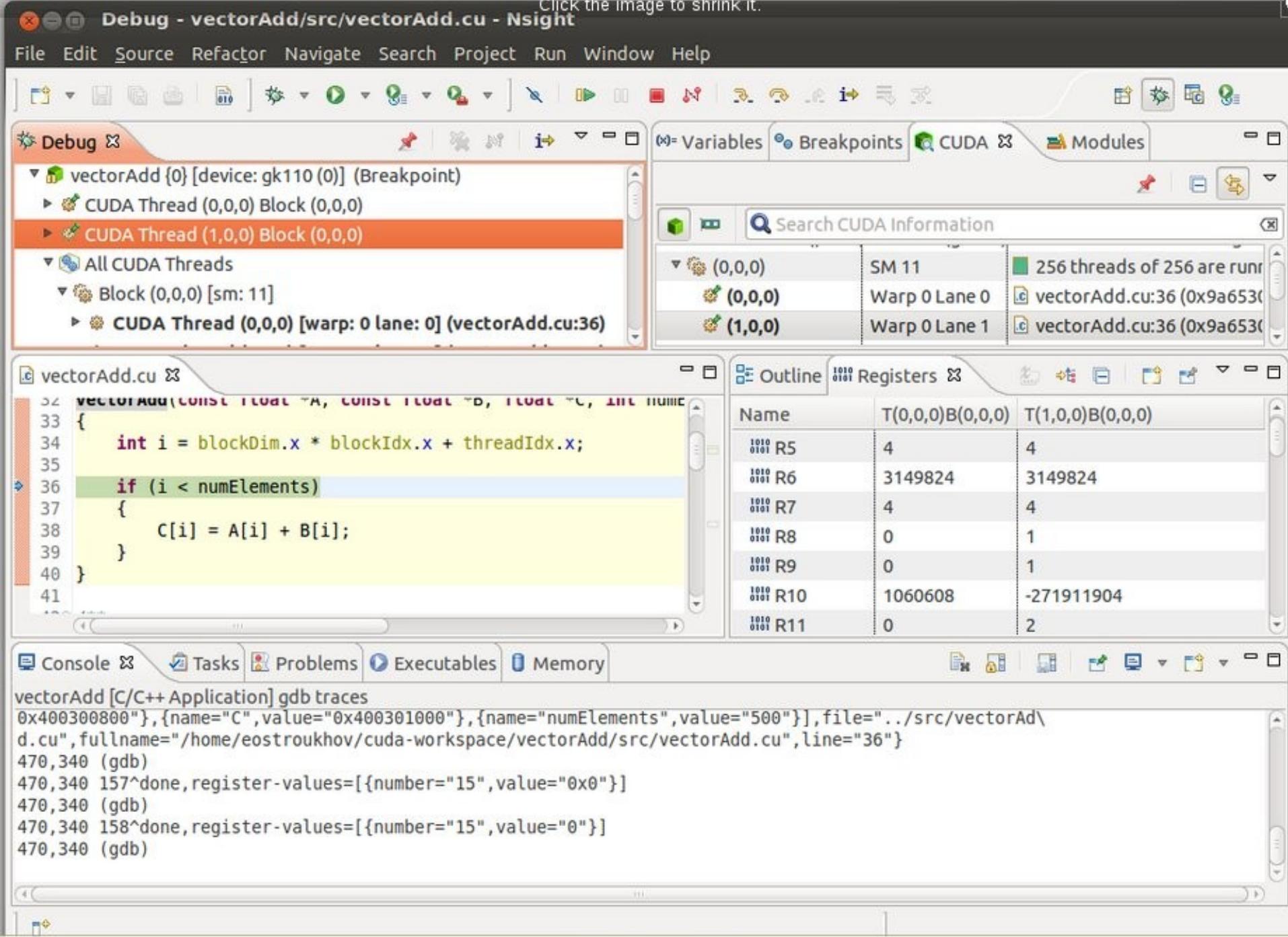
CUDA-GDB runs on Linux and Mac OS X, 32-bit and 64-bit. CUDA-GDB is based on GDB 7.6 on both Linux and Mac OS X.

1.2. Supported Features

CUDA-GDB is designed to present the user with a seamless debugging environment that allows simultaneous debugging of both GPU and CPU code within the same application. Just as programming in CUDA C is an extension to C programming, debugging with CUDA-GDB is a natural extension to debugging with GDB. The existing GDB debugging features are inherently present for debugging the host code, and additional features have been provided to support debugging CUDA device code.

CUDA-GDB supports debugging C/C++ and Fortran CUDA applications. (Fortran debugging support is limited to 64-bit Linux operating system) All the C++ features supported by the NVCC compiler can be debugged by CUDA-GDB.

CUDA-GDB allows the user to set breakpoints, to single-step CUDA applications, and also to inspect and modify the memory and variables of any given thread running on the hardware.



Wrapping Up 1

Exercises (CUDA Lectures in afternoon)

- 0. hello, device- first kernel call, hello world, GPU properties
- 1. add - vector addition using one thread in one block only
- 2. add-index - vector addition using blocks in parallel,
one thread per block only.
- 3. add-parallel - vector addition using all blocks and threads in parallel
- 4. dot - scalar product using shared memory of one block
only for reduction
- 5. dot-full - scalar product using shared memory and
atomic add across blocks
- 6. dot-perfect - scalar product; fat threads and final reduction on host.
- 8. histo** - histogram using fat threads and atomic add
on shared and global memory, timing
- 7. matmul** - matrix multiplication with tiled access shared memory
(expect Friday)

Wrapping Up 2

Elements of CUDA C learnt:

threadId.x , blockDim.x, blockDim.y, gridDim.x
(threadId.y, blockDim.y, blockDim.z, gridDim.y)
kernel<<<n,m>>> (...)
kernel<<<dimBlock, dimGrid>>>(...)

device code

__global__
__shared__

cudaMalloc / cudaFree
cudaMemcpy / cudaMemset
cudaGetDeviceProperties
cudaEventCreate, cudaEventRecord,
cudaEventSynchronize, cudaEventElapsedTime,
cudaEventDestroy
AtomicAdd

Threads, Blocks
(matmul coming with 2D grids)
kernel calls
dim3 variable type (matmul)

shared memory on GPU
manage global memory of GPU
copy/set to or from memory
get device properties in program

CUDA profiling
atomic functions

Wrapping Up 3

What we have not yet learnt...

constant
device

Intrinsic Functions (device type)

https://docs.nvidia.com/cuda/cuda-math-api/group__CUDA__MATH__SINGLE.html#group__CUDA__MATH__SINGLE

constant memory on GPU
functions device to device

host

More atomic functions

cudaBindTexture

fat threads for 2D and 3D stencils

cudaStreamCreate, cudaStreamDestroy

using Tensor Cores

functions host to host

using texture memory
thread coalescence opt.
working with CUDA streams

...

Histogram

Chapter in Book of Jason Sanders

<https://wwwstaff.ari.uni-heidelberg.de/spurzem/lehre/WS20/cuda/files/cuda-histograms.pdf>

[Link on our webpage](#)

On kepler: 8_histo

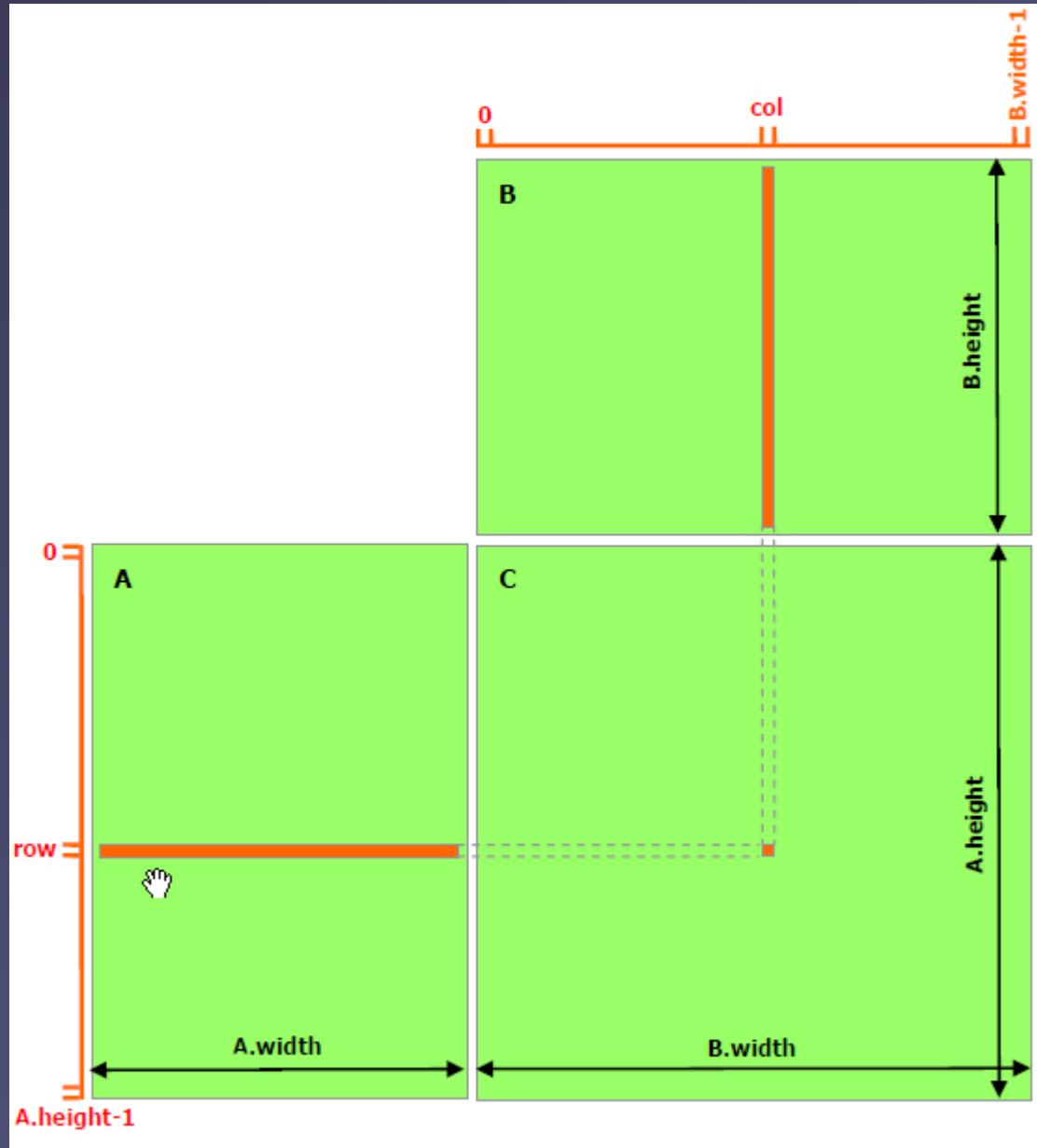
histo.cu

histo-no-atomic.cu

Both use atomic on shared memory!

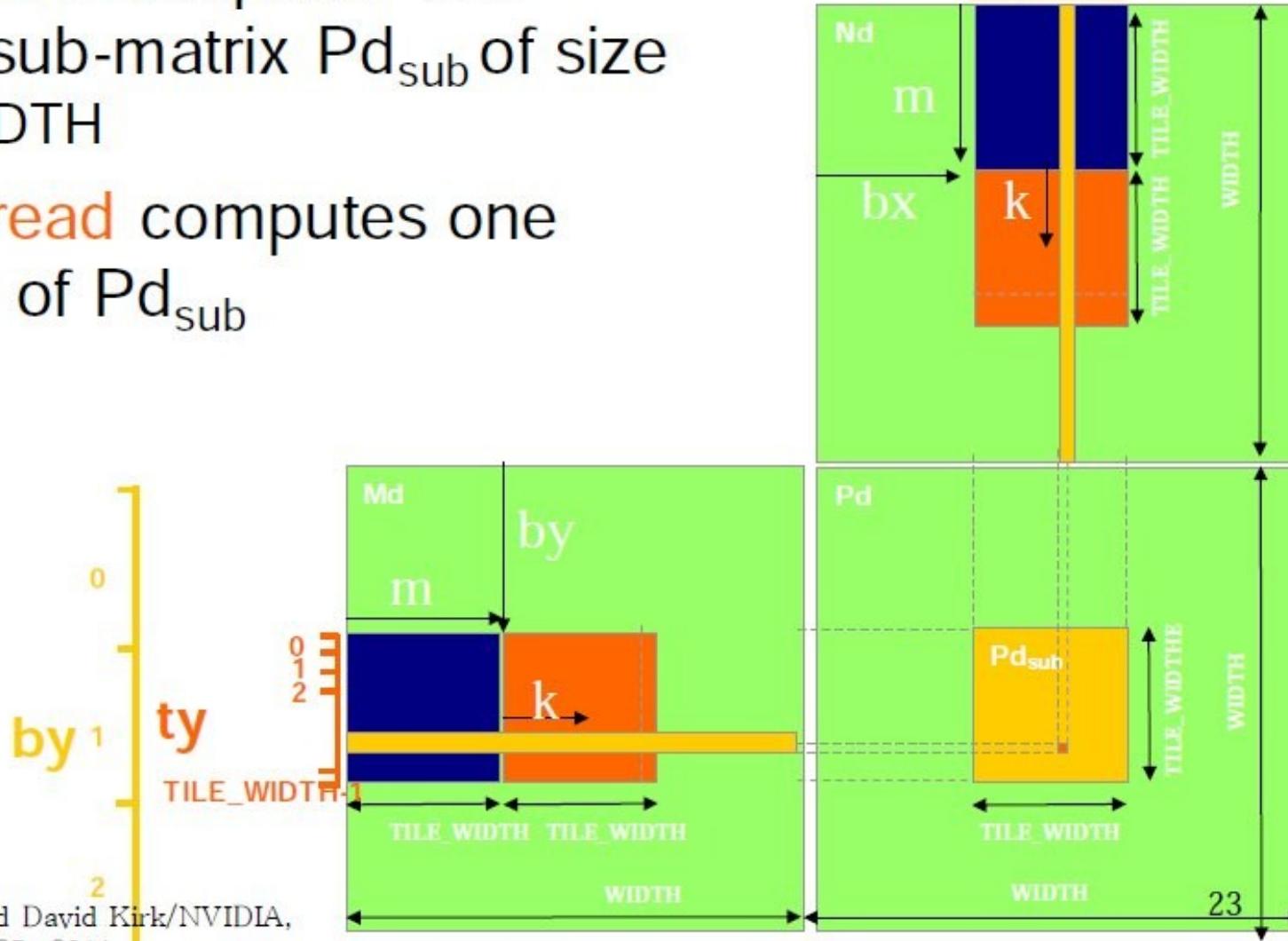
But only first one uses also atomic on global memory!

Intuitive multiply



Tiled Multiply

- Each **block** computes one square sub-matrix Pd_{sub} of size $TILE_WIDTH$
- Each **thread** computes one element of Pd_{sub}



Speed-Up Ratio

GPU speed-up over CPU

